

Explosive appearance of cores and bootstrap percolation on lattices

Lyuben Lichev, IST Austria

k -core of G : its largest subgraph with minimum degree k .

Easy to construct: iteratively delete vertices of smaller degree.

For $k \geq 3$, Łuczak showed in 1991 that the appearance of a k -core in the random graphs process undergoes a discontinuous phase transition.

k -cores of random graphs are typically k -connected.

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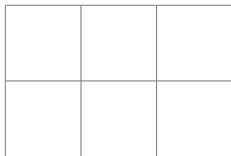
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Our process

Ambient space: $n \times n$ torus in



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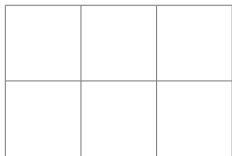


Dynamic random initial condition: vertices arrive consecutively in a random order.

Object of interest: k -core of the graph induced by the present vertices.

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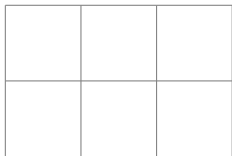


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Examples

On the square lattice:

- the 2-core emerges as a union of (possibly intersecting) short cycles,
- the 4-core emerges when all vertices appeared.

What about the 3-core?

Theorem (Bollobás, Duminil-Copin, Morris, Smith, 2014)

For large $C > 0$ and small $c > 0$, w.h.p.

- *at time $n^2 - Cn^2 / \log n$, the 3-core is empty.*
- *at time $n^2 - cn^2 / \log n$, the 3-core contains $n^2 - o(n^2)$ vertices.*

Theorem (Hartarsky, L., 2024++)

W.h.p. the size of the 3-core jumps from 0 to $n^2 - o(n^2)$ in a single step.

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Connection to bootstrap percolation

State space: $\{\circ, \bullet\}^{n \times n}$ (\circ/\bullet = healthy/infected)

Vertex x in state \circ with at least 2 neighbours \bullet : becomes \bullet

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On the square lattice: 3-core process \iff 2-neighbour bootstrap percolation in the complement

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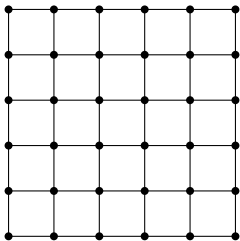
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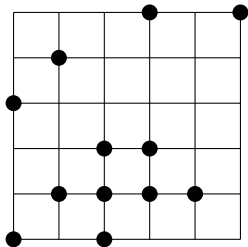
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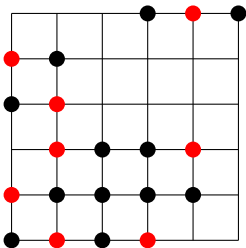
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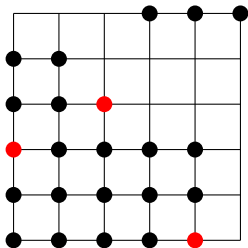
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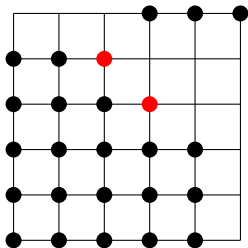
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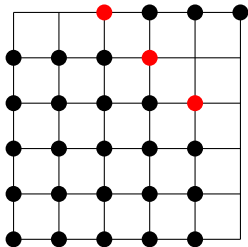
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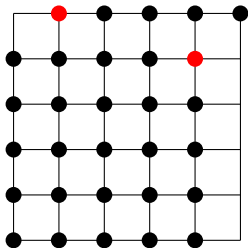
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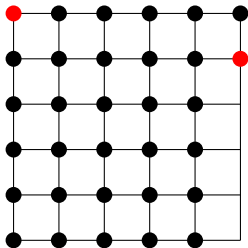
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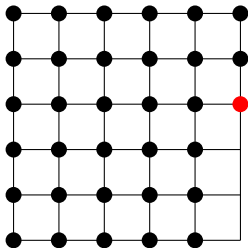
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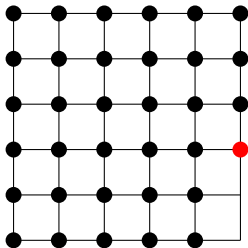
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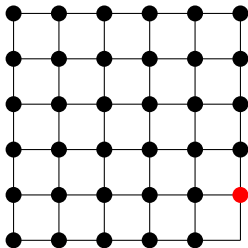
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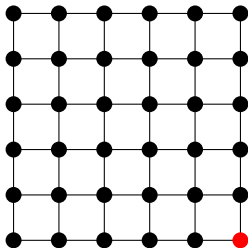
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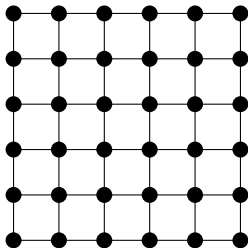
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Why is 2-neighbour bootstrap special?

On the square lattice:

- a single • does infect everyone in 1-neighbour bootstrap (supercritical).
- many finite sets of • remain • forever in 3-neighbour bootstrap (subcritical).

2-neighbour bootstrap interpolates between the two: a critical model.

Theorem (Bollobás, Duminil-Copin, Morris, Smith, 2014)

In 2-neighbour bootstrap, for large $C > 0$ and small $c > 0$, w.h.p.

- $Cn^2 / \log n$ vertices infect the entire torus.
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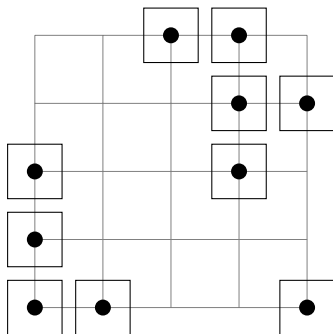
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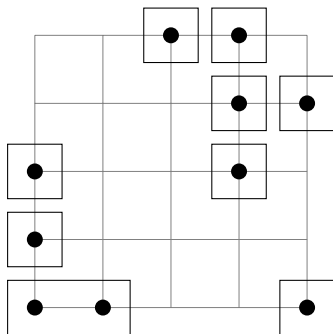
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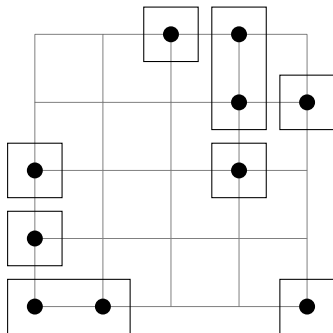
Elements of proof: the rectangles process



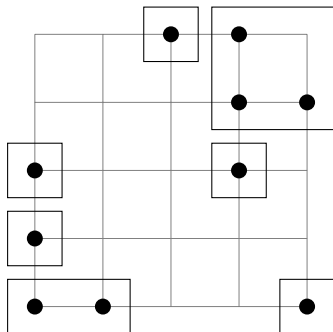
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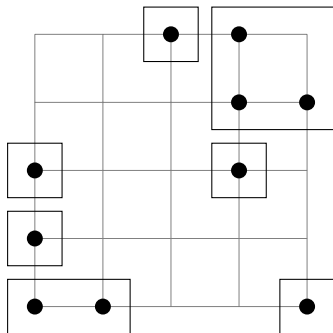
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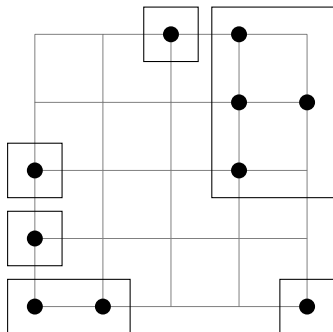
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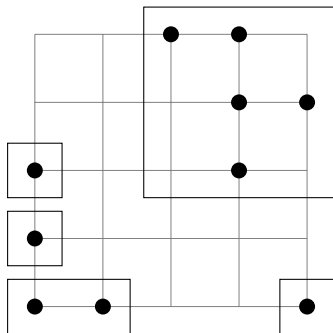
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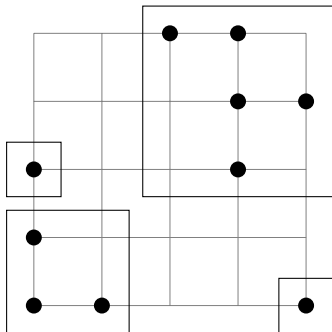
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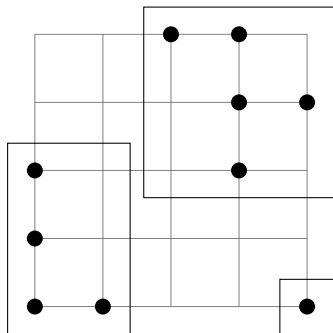
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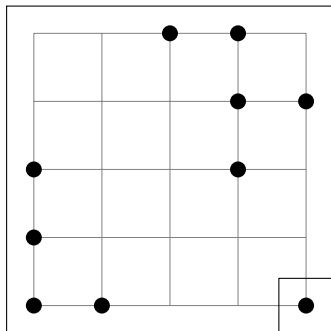
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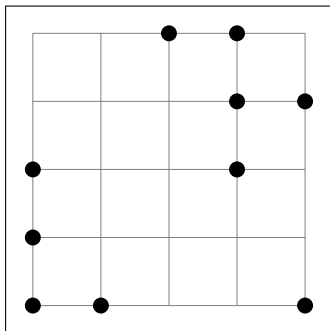
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Elements of proof: using the rectangles process

$\tau =$ smallest time when everyone becomes \bullet

Sharp threshold result: $\tau = \Theta(n^2 / \log n)$ w.h.p.

At time $\tau - 1$, a typical box Λ of size $D(\log n)^3 \times D(\log n)^3$ whose centre is \circ contains only small groups of \bullet -s far from each other.

Conclusion: it stays almost stationary without \bullet -s helping from outside.

If its centre becomes \bullet using \bullet -s outside Λ , then a large infected rectangle is found. Such rectangle would spread and infect the entire torus, impossible at time $\tau - 1$.

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An extension

We manage to extend our hitting-time result to other critical models.

Fix $K \subseteq \mathbb{R}^2$ convex and invariant under $\pi/2$ -rotation, and $\mathcal{K}_s = \mathbb{Z}^2 \cap (sK)$. Choose

$$r = 1 + \min_{u \in \mathbb{R}^2 \setminus \{0\}} \left| \{x \in \mathcal{K}_s : \langle x, u \rangle < 0\} \right|.$$

Theorem (Hartarsky, L., 2024++)

In r -neighbour bootstrap with neighbourhood \mathcal{K}_s for large s , the size of the infected set grows from $O(n^2 / \log n)$ to n^2 in one step.

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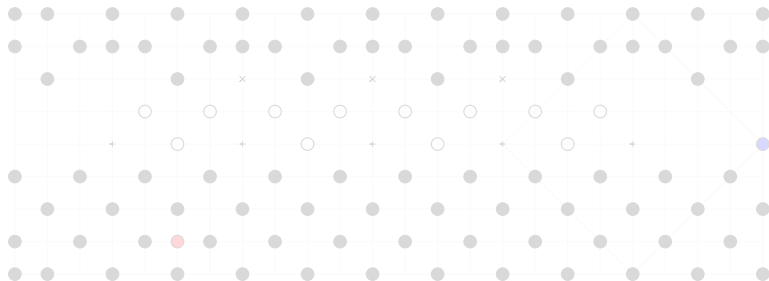
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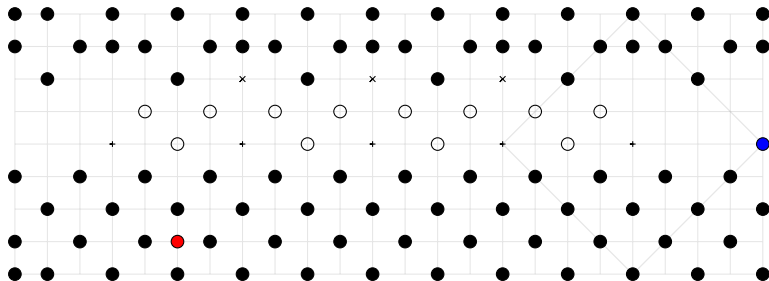
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No rectangle process for large neighbourhoods



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